Responsive Choice in Mobile Processes

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June 12, 2009

¹Joint work with António Ravara

Goal

- Verifying Calculus Encodings
- Using Choice and Branching to encode Data

- Activeness and Responsiveness
- Choice Types
- Conditional Dependencies
- Sound Type System

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$$[\![\mathsf{True}]\!]_a = ! \, a(tf).\bar{t}$$
 $[\![\mathsf{False}]\!]_a = ! \, a(tf).\bar{f}$

$$\llbracket \mathsf{If} (e) \mathsf{Then} (T) \mathsf{Else} (F) \rrbracket = (\nu \mathsf{bt} f) (\llbracket e \rrbracket_b | \overline{b} \langle \mathsf{t} f \rangle | (t. \llbracket T \rrbracket + f. \llbracket F \rrbracket))$$

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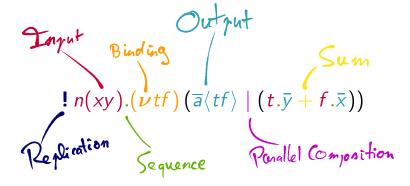
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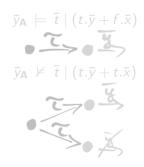
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The Synchronous Polyadic π -calculus

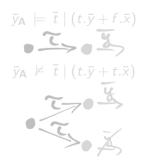


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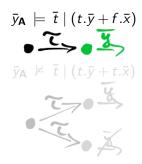
- a_A: I am soon ready to receive at a
- \bar{a}_A : I am soon ready to send to a



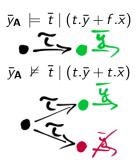
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Definition (R=Responsiveness)

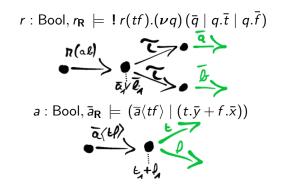
- a_R : If I get a message from you at a, I'll obey a's protocol
- \bar{a}_R : If you get a message from me at a, I'll obey \bar{a} 's protoco

$$r: \mathsf{Bool}, r_{\mathsf{R}} \models ! r(tf).(\nu q) (\bar{q} \mid q.\bar{t} \mid q.\bar{f})$$

$$a: \mathsf{Bool}, \bar{a}_{\mathsf{R}} \models (\bar{a}\langle tf \rangle \mid (t.\bar{y} + f.\bar{x}))$$

Definition (**R**=Responsiveness)

- a_R : If I get a message from you at a, I'll obey a's protocol
- \bar{a}_R : If you get a message from me at a, I'll obey \bar{a} 's protocol



Multiplicities

Definition (Multiplicity)

 p^m : I'll use p at most m times.

- p^0 : Never
- p^1 : Once
- p^{ω} : Once, replicated
- p★: Many times

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Choice

Definition (Choice $A \vee B$)

I will either behave like A or like B

$$\bar{t}_{A} \vee \bar{f}_{A} \models (\nu q) (\bar{q} \mid q.\bar{t} \mid q.\bar{f})$$

Definition (Branching A+B)

You can make me do A or B (not both)

$$(t_{\mathsf{A}}+f_{\mathsf{A}}) \models (t.\bar{y}+f.\bar{x})$$

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Dependencies

Definition (Dependency $A \triangleleft B$)

If you give me B, I'll give you A.

Choice

$$n_{\mathbf{R}} \triangleleft (a_{\mathbf{A}} \wedge a_{\mathbf{R}}) \models ! n(xy).\overline{a}(\nu tf).(t.\overline{y} + f.\overline{x})$$



Types

Definition (Channel Type)

(parameters ; input ; output)

- Bool $\stackrel{\text{def}}{=}$ (()() ; $\bar{1}_{\mathbf{A}} \vee \bar{2}_{\mathbf{A}}$; $1_{\mathbf{A}} + 2_{\mathbf{A}}$)
- ! $a(xy).(\nu q)(\bar{q}|q.\bar{x}|q.\bar{y}) | \bar{a}\langle tf \rangle | (t.P+f.Q)$

Choice

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Typing — Composition

$$(\bar{t}_{\mathsf{A}} \vee \bar{f}_{\mathsf{A}}) \triangleleft \mathsf{a}_{\mathsf{AR}} \vdash \bar{\mathsf{a}} \langle tf \rangle$$

$$((t_{\mathsf{A}}, \bar{y}_{\mathsf{A}} \triangleleft \bar{t}_{\mathsf{A}}) + (f_{\mathsf{A}}, \bar{x}_{\mathsf{A}} \triangleleft \bar{f}_{\mathsf{A}})) \vdash (t.\bar{y} + f.\bar{x})$$

$$(\bar{t}_{\mathsf{A}} \vee \bar{f}_{\mathsf{A}}) \triangleleft \mathsf{a}_{\mathsf{AR}}, ((t_{\mathsf{A}}, \bar{y}_{\mathsf{A}} \triangleleft \bar{t}_{\mathsf{A}}) + (f_{\mathsf{A}}, \bar{x}_{\mathsf{A}} \triangleleft \bar{f}_{\mathsf{A}})) \vdash \bar{\mathsf{a}} \langle tf \rangle \mid (t.\bar{y} + f.\bar{x})$$

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 $(\bar{t}_{\mathbf{A}} \vee \bar{f}_{\mathbf{A}}) \triangleleft a_{\mathbf{A}\mathbf{R}} \vdash \bar{a} \langle tf \rangle$

Typing — Composition

$$(\overline{t}_{\mathbf{A}} \vee \overline{f}_{\mathbf{A}}) \triangleleft a_{\mathbf{A}\mathbf{R}} \vdash \overline{a} \langle tf \rangle$$

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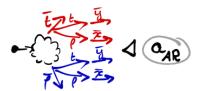
$$(\overline{t}_{\mathbf{A}} \vee \overline{f}_{\mathbf{A}}) \triangleleft a_{\mathbf{A}\mathbf{R}}, ((t_{\mathbf{A}}, \overline{y}_{\mathbf{A}} \triangleleft \overline{t}_{\mathbf{A}}) + (f_{\mathbf{A}}, \overline{x}_{\mathbf{A}} \triangleleft \overline{f}_{\mathbf{A}})) \vdash \overline{a} \langle tf \rangle \mid (t.\overline{y} + f.\overline{x})$$

Typing — ∨-Distributivity

∨-Distributivity

"V distributes with everything else"

$$(\overline{t}_{\mathsf{A}} \lor \overline{f}_{\mathsf{A}}) \triangleleft a_{\mathsf{AR}}, ((t_{\mathsf{A}}, \overline{y}_{\mathsf{A}} \triangleleft \overline{t}_{\mathsf{A}}) + (f_{\mathsf{A}}, \overline{x}_{\mathsf{A}} \triangleleft \overline{f}_{\mathsf{A}})) =$$



Branching Collapse

$$(\bar{p}_{\mathbf{A}}, (p_{\mathbf{A}} + \varepsilon)) \hookrightarrow (\bar{p}_{\mathbf{A}} \wedge p_{\mathbf{A}}) \vee (\bar{p}_{\mathbf{A}}, (p_{\mathbf{A}} + \varepsilon))$$

$$(\bar{t}_{\mathsf{A}}, t_{\mathsf{A}}, (\bar{y}_{\mathsf{A}} \triangleleft \bar{t}_{\mathsf{A}})) \triangleleft a_{\mathsf{A}\mathsf{R}} \quad \lor \quad (\bar{f}_{\mathsf{A}}, f_{\mathsf{A}}, (\bar{x}_{\mathsf{A}} \triangleleft \bar{f}_{\mathsf{A}})) \triangleleft a_{\mathsf{A}\mathsf{R}}$$



Branching Collapse

$$(\bar{p}_{A}, (p_{A} + \varepsilon)) \hookrightarrow (\bar{p}_{A} \wedge p_{A}) \vee (\bar{p}_{A}, (p_{A} + \varepsilon))$$

$$(\alpha \triangleleft \beta) \land (\beta \triangleleft \gamma) \quad \hookrightarrow \quad (\alpha \triangleleft \beta \lor \gamma) \land (\beta \triangleleft \gamma)$$

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$$((t_{A} \wedge \overline{t}_{A}) \triangleleft a_{AR}, (\overline{y}_{A} \triangleleft \overline{t}_{A})) \vee ((f_{A} \wedge \overline{f}_{A}) \triangleleft a_{AR}, (\overline{x}_{A} \triangleleft \overline{f}_{A})) \hookrightarrow ((t_{A}, \overline{t}_{A}, \overline{y}_{A}) \triangleleft a_{AR}) \vee ((f_{A}, \overline{x}_{A}, \overline{f}_{A}) \triangleleft a_{AR}) =$$

$$((t_{A}, \overline{t}_{A}, \overline{y}_{A}) \vee (f_{A}, \overline{x}_{A}, \overline{f}_{A})) \triangleleft a_{AR} \vdash \overline{a}(tf) \mid (t, \overline{y} + f, \overline{x})$$

$$(\alpha \triangleleft \beta) \land (\beta \triangleleft \gamma) \quad \hookrightarrow \quad (\alpha \triangleleft \beta \lor \gamma) \land (\beta \triangleleft \gamma)$$

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Typing — Responsiveness

$$n_{\mathbf{R}} \triangleleft \ldots \vdash P = ! n(xy).(\nu tf) (\overline{a}\langle tf \rangle \mid (t.\overline{y} + f.\overline{x}))$$

Responsiveness

$$(\nu \tilde{y}) (n : \sigma \odot n_{\mathbb{R}} \triangleleft \sigma \{\tilde{y}/_{1\cdots n}\} \odot \cdots) \vdash n(\tilde{y})$$

$$(\nu xy) \left(n_{R} \triangleleft (\bar{x}_{A} \vee \bar{y}_{A}) \wedge (\bar{x}_{A} \vee \bar{y}_{A}) \triangleleft a_{AR} \right) \vdash P$$

$$n_{R} \triangleleft \bullet P$$

$$\uparrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad$$

Typing — Responsiveness

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Responsiveness

$$(\boldsymbol{\nu}\tilde{\mathbf{y}})(\mathbf{n}:\sigma\odot n_{\mathbf{R}}\triangleleft\sigma\{\tilde{\mathbf{y}}/_{1\cdots n}\}\odot\cdots)\vdash n(\tilde{\mathbf{y}})$$

$$\hookrightarrow n_{\mathbf{R}} \triangleleft a_{\mathbf{A}\mathbf{R}} \vdash ! n(xy).(\nu tf) (\overline{a}\langle tf \rangle \mid (t.\overline{y} + f.\overline{x}))$$





Typing — Responsiveness

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Responsiveness

$$(\boldsymbol{\nu}\tilde{\mathbf{y}}) (n : \sigma \odot n_{\mathbf{R}} \triangleleft \sigma \{\tilde{\mathbf{y}}/_{1\cdots n}\} \odot \cdots) \vdash n(\tilde{\mathbf{y}})$$

Conclusion

Our contribution:

- ullet A formalism describing liveness properties in the π -calculus
- Choice and Branching types ⇒ Data encodings
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Thank you!



More info:

• http://maxime.gamboni.org/